

# Somatorios

# Arithmetic Series

$$n + (n-1) + (n-2) + \dots + 2 + 1 = n(n+1)/2$$

Typical case:

For i = 1 to n

For j = i+1 to n

*do something*

End

End

For a fixed value of i, the inner "For" does  $n-(i+1)+1 = n-i$  iterations

Iterations of inner  
"For" when i = 1

Iterations of inner  
"For" when i = 2

Number of iterations:  $(n-1) + (n-2) + \dots + 2 + 1 = (n-1)n/2$

# Geometric Series

- $1 + \frac{1}{2} + \frac{1}{4} + \dots + \frac{1}{2^i} + \dots = 2$
- More generally for  $0 \leq r < 1$ ,  $1 + r + r^2 + \dots = \frac{1}{1-r}$
- A finite Geometric Series can be upper bounded by the infinite one

Ex:  $1 + \frac{1}{2} + \frac{1}{4} + \dots + \frac{1}{2^n} \leq 1 + \frac{1}{2} + \frac{1}{4} \dots = 2$

Obs: Usually  $1/2^i$  becomes small so quickly that even  $\sum_i b_i/2^i$  is at most a constant, for most  $b_i$ 's we will see

Ex:  $\sum_i i/2^i$  is at most a constant  
 $\sum_i i^2/2^i$  is at most a constant

# General upper bound

- Upper bound each term by the maximum
- $a_1 + \dots + a_n \leq n \max \{a_i\}$

Ex: Show that  $1^2 + 2^2 + \dots + n^2 = O(n^3)$

- This is the bound we will use the most
- But it is bad in some cases, for instance for Geometric Series  
 $1 + \frac{1}{2} + \frac{1}{4} + \dots + \frac{1}{2^n}$ 
  - This sum is  $\leq 2$ ...
  - ... but upper bounding each term gives  $\leq n$

# General lower bound

- Want to show  $a_1 + a_2 + \dots + a_n \geq \text{value}$
- Can lower bound each term by minimum  
 $a_1 + \dots + a_n \geq n \min \{a_i\}$
- But this is usually very bad: look at  $n + (n-1) + (n-2) + \dots + 2 + 1$   
 $= n(n+1)/2$ , which is  $\Omega(n^2)$ ...  
...but lower bounding each term gives  $\geq n * 1 = n$

**Trick:** Discard smaller terms (usually the smallest  $n/2$ )

- Back to example  $n + (n-1) + (n-2) + \dots + 2 + 1$   
 $\geq n + (n-1) + (n-2) + \dots + n/2$  ← Discarded smallest  $n/2$  terms  
 $\geq (n/2) * (n/2) = n^2/4$  ← Lower bounded each term by minimum

# Exercise

Exercise 1: Show that  $1^2+2^2+\dots+n^2 = \theta(n^3)$

# Exercise

Exercise 2: Show that  $\log(n!)$  is  $\Theta(n \log n)$

Answer: First we show that  $\log(n!)$  is  $O(n \log n)$ :

$$\log(n!) = \log n + \log(n-1) + \dots + \log 1 < n \cdot \log n$$

Now we show that  $\log(n!)$  is  $\Omega(n \log n)$ :

$$\begin{aligned} \log(n!) &= \log n + \log(n-1) + \dots + \log 1 \\ &\geq n/2 \cdot \log(n/2) \\ &= n/2 (\log n - 1) \end{aligned}$$